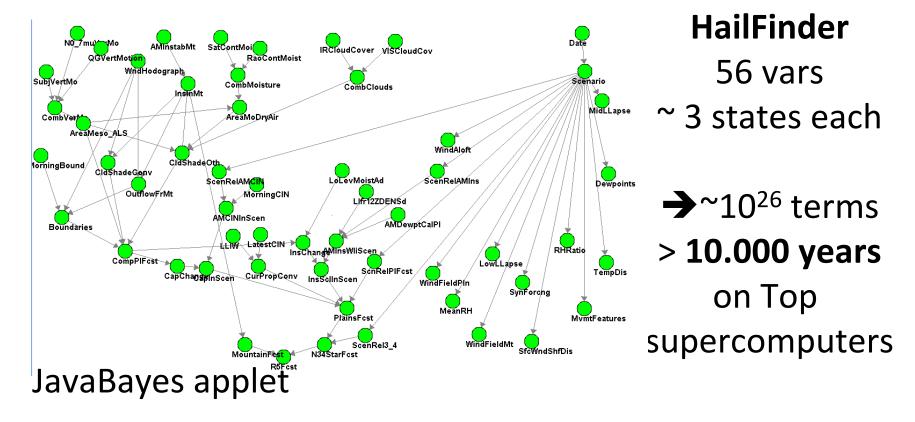
# Probabilistic Graphical Models

#### Lecture 3 – Bayesian Networks Semantics

CS/CNS/EE 155 Andreas Krause

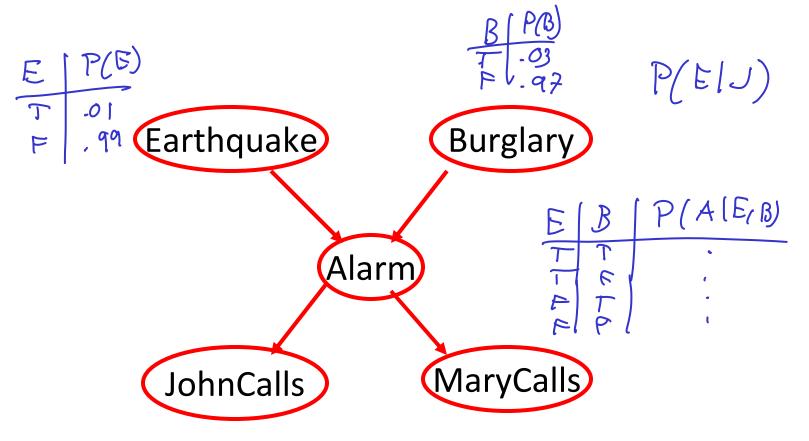
## Bayesian networks

- Compact representation of distributions over large number of variables
- (Often) allows efficient exact inference (computing marginals, etc.)



## Causal parametrization

 Graph with directed edges from (immediate) causes to (immediate) effects



## Bayesian networks

A Bayesian network structure is a directed, acyclic graph G, where each vertex s of G is interpreted as a random variable X<sub>s</sub> (with unspecified distribution)

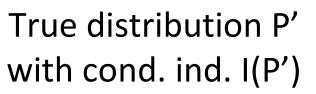
- A Bayesian network (G,P) consists of
  - A BN structure G and ..
  - ..a set of conditional probability distributions (CPDs)
    P(X<sub>s</sub> | Pa<sub>x<sub>s</sub></sub>), where Pa<sub>x<sub>s</sub></sub> are the parents of node X<sub>s</sub> such that
  - (G,P) defines joint distribution

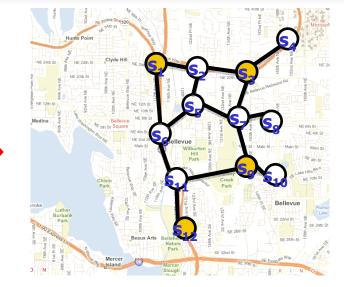
$$P(X_1, ..., X_n) = \prod_i P(X_i \mid \mathbf{Pa}_{X_i})$$

#### Representing the world using BNs

represent





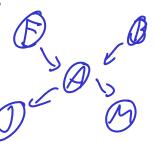


#### Bayes net (G,P) with I(P)

- Want to make sure that  $I(\underline{P}) \subseteq I(\underline{P'})$
- Need to understand CI properties of BN (G,P)

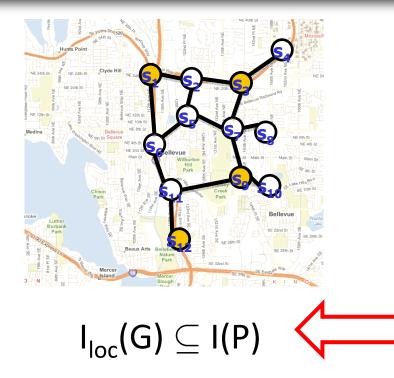
## Local Markov Assumption

- Each BN Structure G is associated with the following conditional independence assumptions
- $J \perp B \mid A$ X  $\perp$  NonDescendents<sub>x</sub> | Pa<sub>x</sub>



- We write I<sub>loc</sub>(G) for these conditional independences
- Suppose (G,P) is a Bayesian network representing P
  Does it hold that I<sub>loc</sub>(G) ⊆ I(P)?
  If this holds, we say G is an I-map for P.

#### **Factorization Theorem**





True distribution P can be represented exactly as Bayesian network (G,P)  $P(X_1,...,X_n) = \prod P(X_i | \mathbf{Pa}_{X_i})$ 

G is an **I-map** of P (independence map)

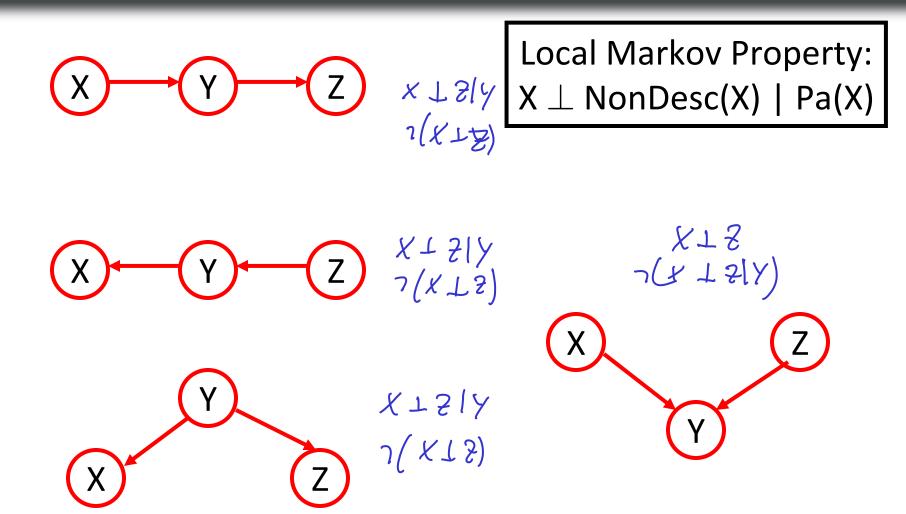
#### Additional conditional independencies

- BN specifies joint distribution through conditional parameterization that satisfies Local Markov Property I<sub>loc</sub>(G) = {(X<sub>i</sub> ⊥ Nondescendants<sub>Xi</sub> | Pa<sub>Xi</sub>)}
- But we also talked about additional properties of CI
  - Weak Union, Intersection, Contraction, ...
- Which additional CI does a particular BN specify?
  - All CI that can be derived through algebraic operations

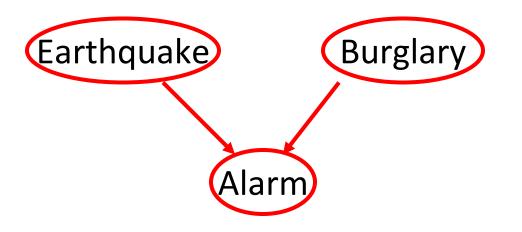
proving Cl is very cumbersome!!

Is there an easy way to find all independences of a BN just by looking at its graph??



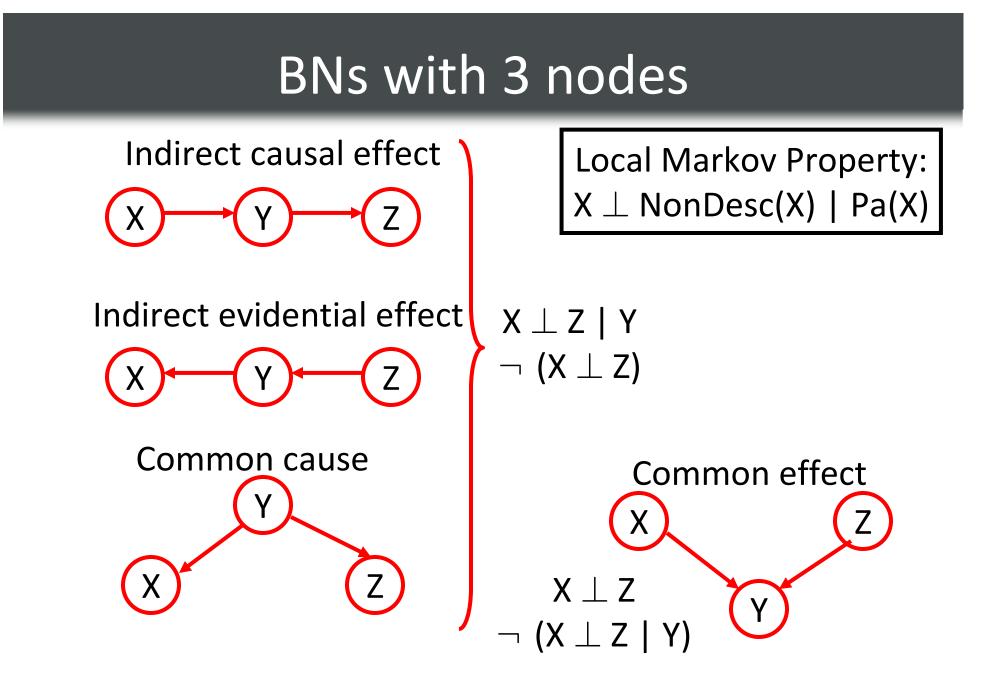


#### V-structures

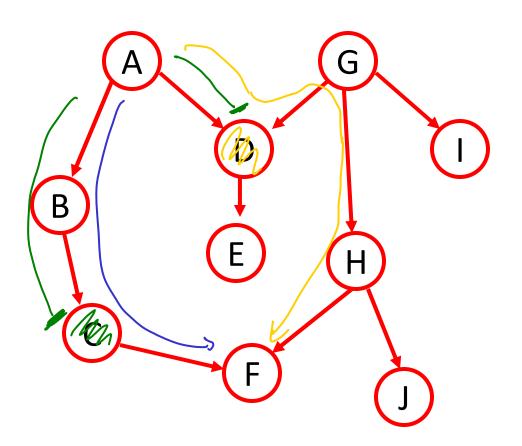


• Know  $E \perp B$ 

• Suppose we know A. Does  $E \perp B \mid A \text{ hold}$ ? (a Loppon:  $P(E = T \mid A = T, B = T) < P(E = T \mid A = T)$ Explaining away

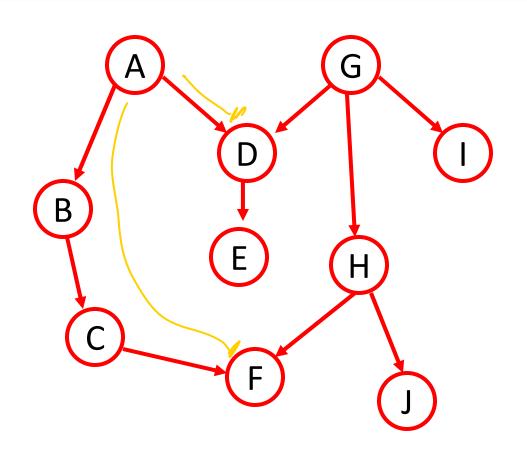


### Examples



 $A \perp F$   $A \perp F \mid C$  $A \perp P \mid C \mid X$ 

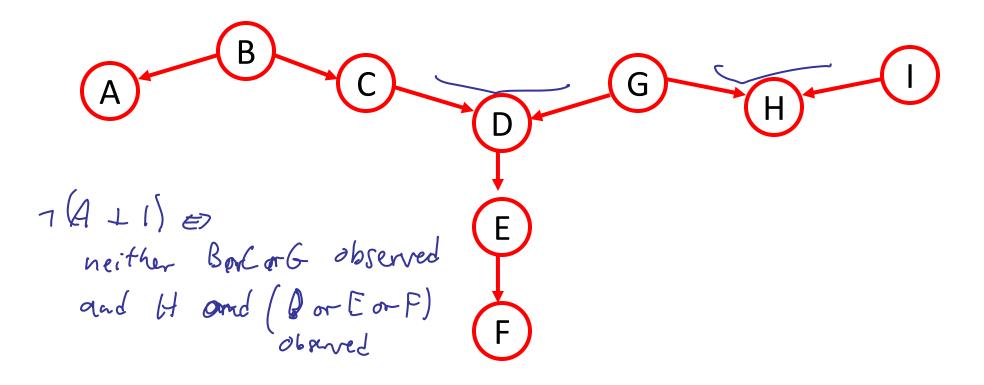
#### More examples



 $\begin{array}{c} A \perp G \\ A \perp G \mid D \times \\ A \perp G \mid E \end{array}$ 

#### Active trails

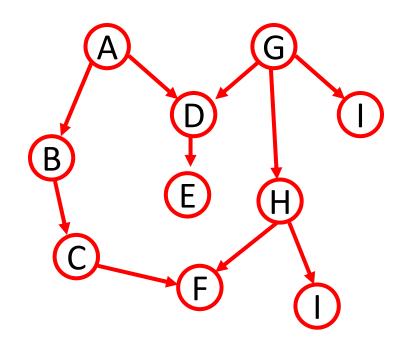
When are A and I independent?



## Active trails

- An undirected path in BN structure G is called active trail for observed variables  $O \subseteq \{X_1, ..., X_n\}$ , if for every consecutive triple of vars X,Y,Z on the path
- Vlock  $X \rightarrow Y \rightarrow Z$  and Y is unobserved  $(Y \notin \mathbf{O})$   $Y \in O$   $X \leftarrow Y \leftarrow Z$  and Y is unobserved  $(Y \notin \mathbf{O})$   $X \leftarrow Y \rightarrow Z$  and Y is unobserved  $(Y \notin \mathbf{O})$ 
  - - X → Y ← Z and Y or any of Y's descendants is observed
    - Any variables  $X_i$  and  $X_i$  for which  $\nexists$  active trail for observations O are called d-separated by O We write **d-sep(X<sub>i</sub>;X<sub>i</sub> | O)**
    - Sets A and B are d-separated given O if d-sep(X,Y | O) for all X ext{A}, Y ext{B}. Write d-sep(A; B | O)

#### d-separation and independence



Theorem: d-sep(X;Y | Z)  $\rightarrow$  X  $\perp$  Y | Z

i.e., X cond. ind. Y given Z if there does not exist any active trail between X and Y for observations **Z** 

 Proof uses algebraic properties of conditional independence

## Soundness of d-separation

• Have seen: P factorizes according to  $G \Leftrightarrow I_{loc}(G) \subseteq I(P)$ 

• Define I(G) = {(X  $\perp$  Y | Z): d-sep<sub>G</sub>(X;Y |Z)}

Theorem: Soundness of d-separation
 P factorizes over G → I(G) ⊆ I(P)

- Hence, d-separation captures only true independences
- How about I(G) = I(P)?

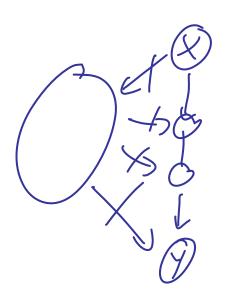
## Does the converse hold?

Suppose P factorizes over G. Does it hold that  $I(P) \subseteq I(G)$ ?  $P \vdash X \perp Y \qquad I(P) = \{(X \perp Y)\}$  $G: \otimes \neg \otimes \qquad I(G) \geq \{\}$ 

#### Existence of dependences for non-d-separated variables

 Theorem: If X and Y are not d-separated given Z, then there exists some distribution P factorizing over G in which X and Y are dependent given Z

Proof sketch:



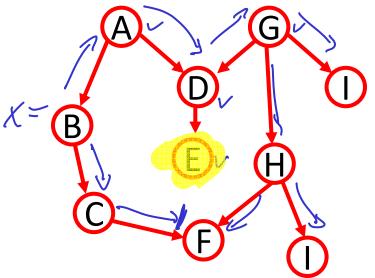
Pick active trail Parameterize CDDs along trail to create dependence Everything close set to independent to avoid concelling dependencies

#### Completeness of d-separation

- Theorem: For "almost all" distributions P that factorize over G it holds that I(G) = I(P)
  - "almost all": except for a set of distributions with measure 0, assuming only that no finite set of distributions has measure > 0

## Algorithm for d-separation

- How can we check if  $X \perp Y \mid Z$ ?
  - Idea: Check every possible path connecting X and Y and verify conditions
  - Exponentially many paths!!! 🙁
- Linear time algorithm:
  Find all nodes reachable from X
  - 1. Mark **Z** and its ancestors
  - 2. Do breadth-first search starting from X; stop if path is blocked
  - Have to be careful with implementation details (see reading)

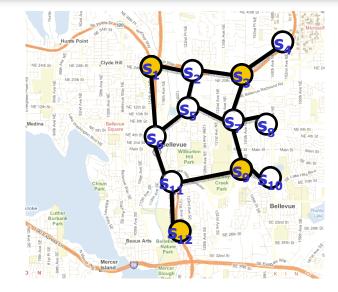


#### Representing the world using BNs

represent



True distribution P' with cond. ind. I(P')



#### Bayes net (G,P) with I(P)

- Want to make sure that  $I(P) \subseteq I(P')$
- Ideally: I(P) = I(P')
- Want BN that exactly captures independencies in P'!

## Minimal I-maps

- Lemma: Suppose G' is derived from G by adding edges
- Then I(G')  $\subseteq$  I(G)
- Proof:

(ompleteness: ICG) = { all CIs derivable from loc(G) using CI properties 3 => I(G') < I(G) J

- Thus, want to find graph G with I(G) ⊆ I(P) such that when we remove any single edge, for the resulting graph G' it holds that I(G') ⊈ I(P)
- Such a graph G is called minimal I-map

 $|_{(G')} \leq |_{Loc} (G)$ 

## Existence of Minimal I-Maps

Does every distribution have a minimal I-Map?

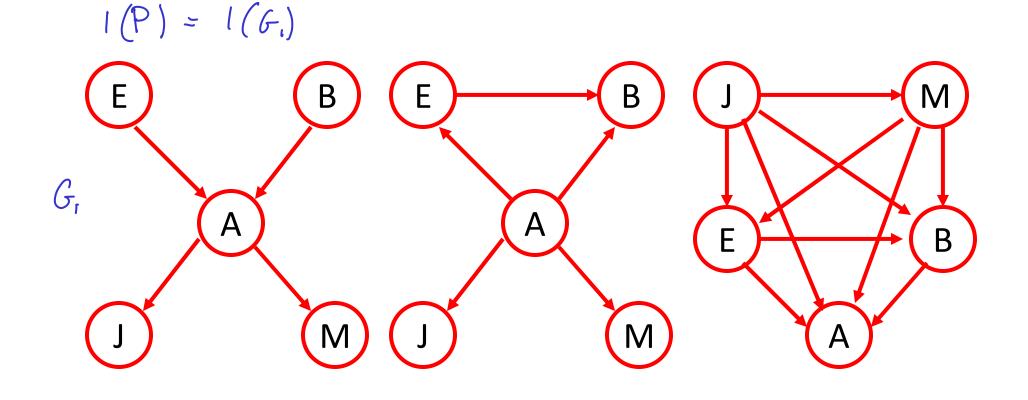
#### Algorithm for finding minimal I-map

- Given random variables and known conditional independences
- Pick ordering X<sub>1</sub>,...,X<sub>n</sub> of the variables
- For each X<sub>i</sub>
  - Find minimal subset  $\mathbf{A} \subseteq \{X_1, ..., X_{i-1}\}$  such that  $P(X_i \mid X_1, ..., X_{i-1}) = P(X_i \mid \mathbf{A})$

#### Will produce minimal I-map!

## Uniqueness of Minimal I-maps

Is the minimal I-Map unique?

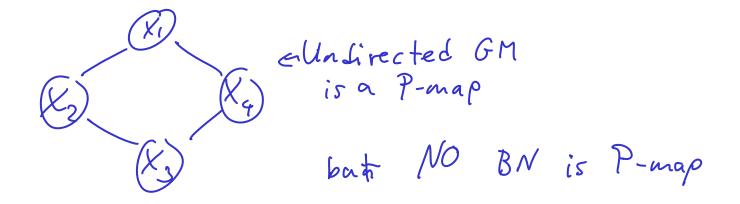


## Perfect maps

- Minimal I-maps are easy to find, but can contain many unnecessary dependencies.
- A BN structure G is called P-map (perfect map) for distribution P if I(G) = I(P)
- Does every distribution P have a P-map?

## Existence of perfect maps

#### Existence of perfect maps

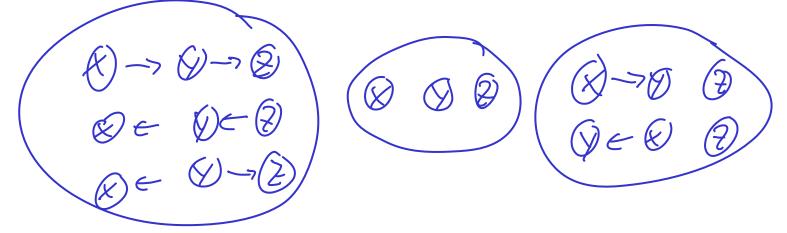


#### Uniqueness of perfect maps

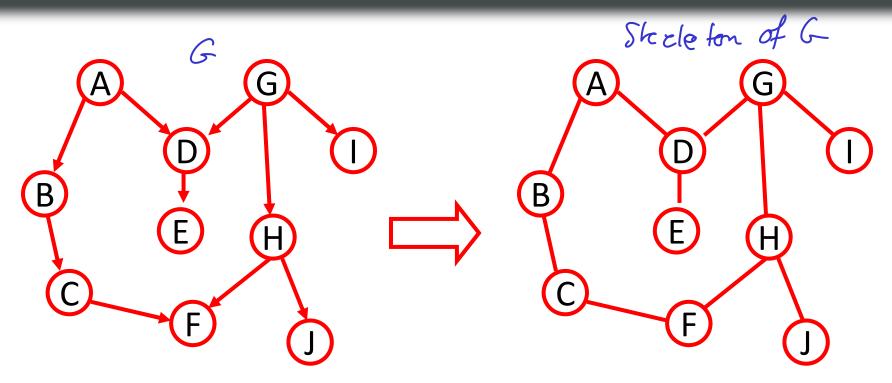
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## I-Equivalence

- Two graphs G, G' are called I-equivalent if I(G) = I(G')
- I-equivalence partitions graphs into equivalence classes



#### **Skeletons of BNs**



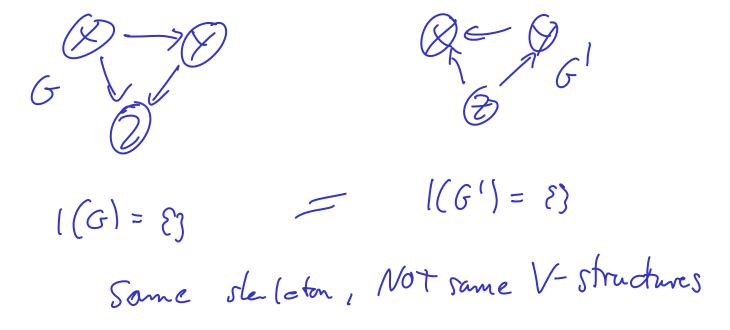
I-equivalent BNs must have same skeleton

Some skeleton \$ 1-equiv. Q. (9)

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#### Importance of V-structures

- Theorem: If G, G' have same skeleton and same Vstructure, then I(G) = I(G')
- Does the converse hold?



## Immoralities and I-equivalence

 A V-structure X → Y ← Z is called immoral if there is no edge between X and Z ("unmarried parents")

immorality

Theorem: I(G) = I(G') ⇔ G and G' have the same skeleton and the same immoralities.

## Tasks

- Subscribe to Mailing list <u>https://utils.its.caltech.edu/mailman/listinfo/cs155</u>
- Read Koller & Friedman Chapter 3.3-3.6
- Form groups and think about class projects. If you have difficulty finding a group, email Pete Trautman
- Homework 1 out tonight, due in 2 weeks. Start early!