Introduction to Artificial Intelligence

Lecture 4 – Adversarial search

CS/CNS/EE 154

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Projects

- Recitations: Thursday 4:30pm 5:30pm, Annenberg 107
 - Details about projects
 - Will also be posted on webpage
- By Monday 10/11
 - Form team of 3 students
 - Need to select project (Doodle link will be sent today)
 - For independent projects: need to submit proposal
- If you don't have a team, send email to TAs
- Homework 1 out on Friday

Types of games

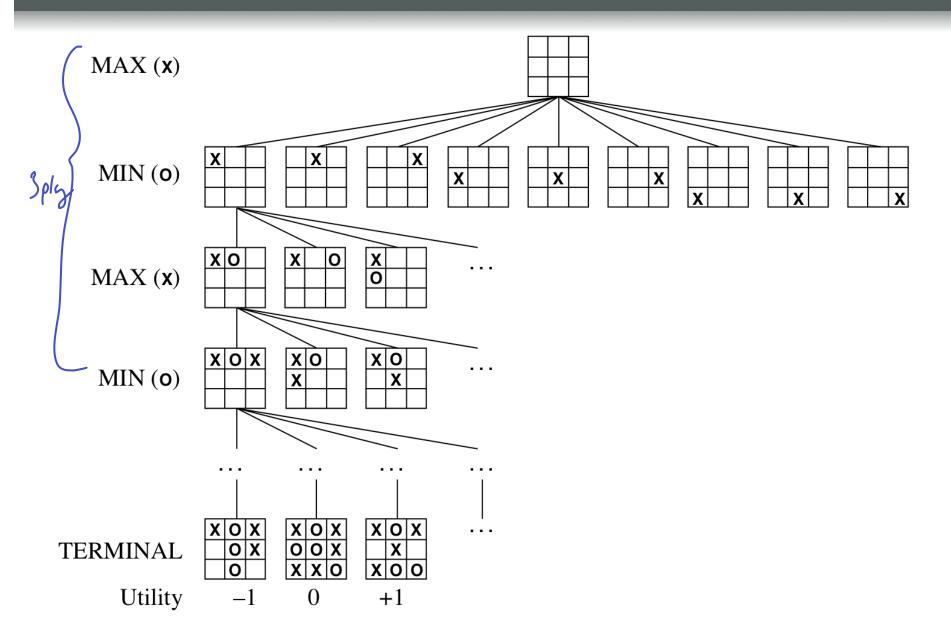
	Chess	Backgammon	Poker	Rock Paper Scissors	WoW
Observable?	Y	Y	\mathcal{N}	> *	N
Determ.?	Y	N	N		N
Simultan.?	N	V	N	Y	У
Zero-sum?	Y	Y	Y	Y	N
Discrete?	Y	Y	Y	Y	N
# Players?	2	2	≥ 2	2	≈ 8·10 ⁶

In this class, focus on two-player, sequential, zero-sum, discrete (mostly deterministic)

Games vs. search

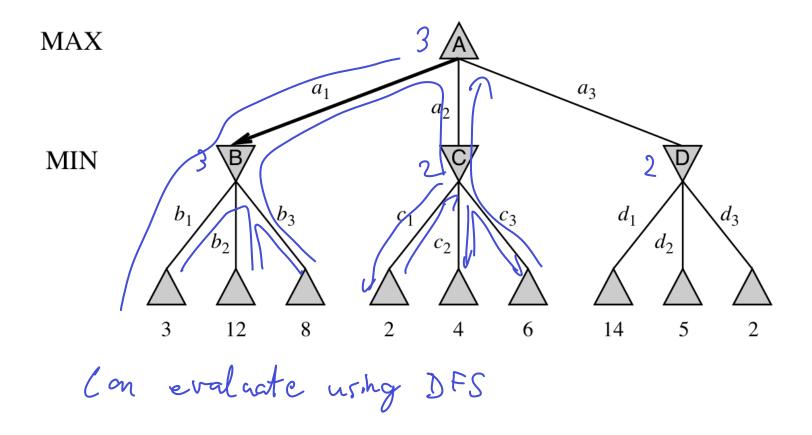
- In games, actions are nondeterministic
 - Opponent can affect state of the environment
- Optimal solution no longer sequence of actions, instead a strategy (policy, conditional plan)
 - If you X I'll do Y, else if you do Y I'll do Z,

Game tree



Minimax game tree

- Search for optimal move no matter what opponent does
- minimax value = best achievable payoff against best play



Solving deterministic games

- MiniMax used to calculate optimal move:
- Inductive definition:

If *n* is terminal node:

Value is utility(n.state)

If *n* is MAX node:

Value is highest value of all successor node values

If *n* is MIN node

Value is lowest value of all successor node values

Properties of minimax search

- · Complete? Yes if tree is finite
- Time complexity?
 O(b^m)
- Space complexity? O(b·m)
- Optimal? Yes if finite and

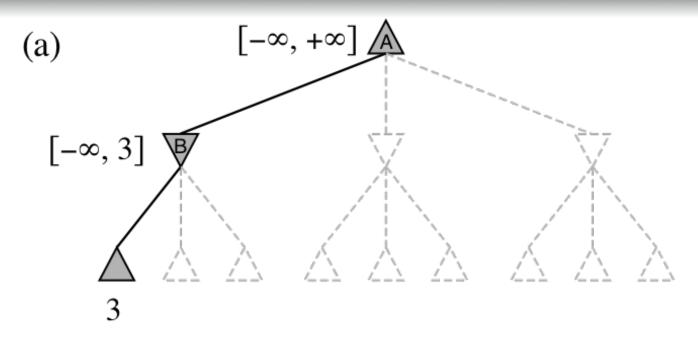
 if opponent plays optimally

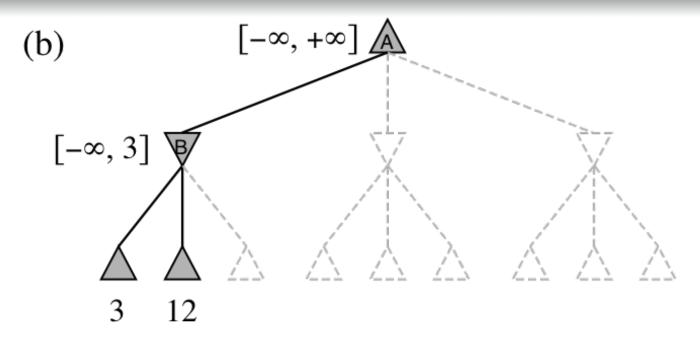
 For Tictactoer. b = 9, m = 9 9; elin symmetrics

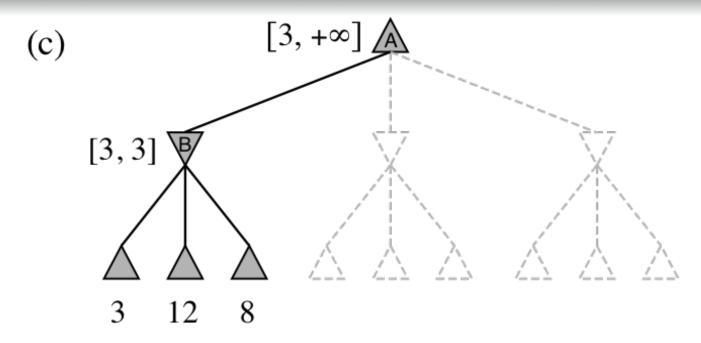
 >> = 30000 nodes

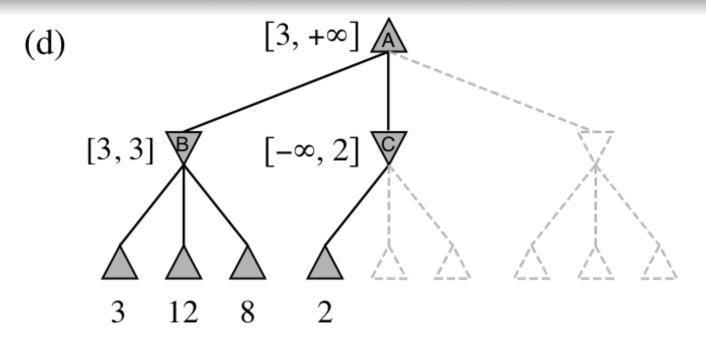
 For Chess: b = 35, m = 50, 35 60 hopelessy intractable;

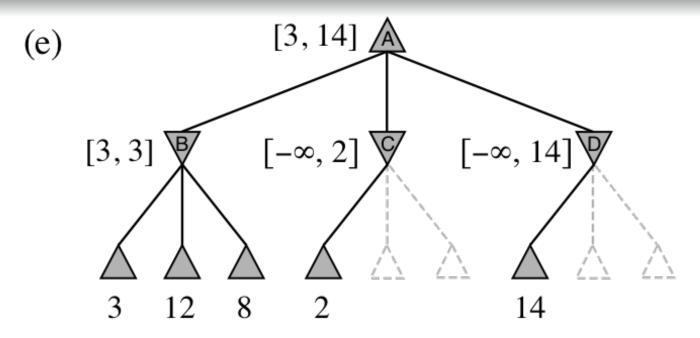
α-β-pruning

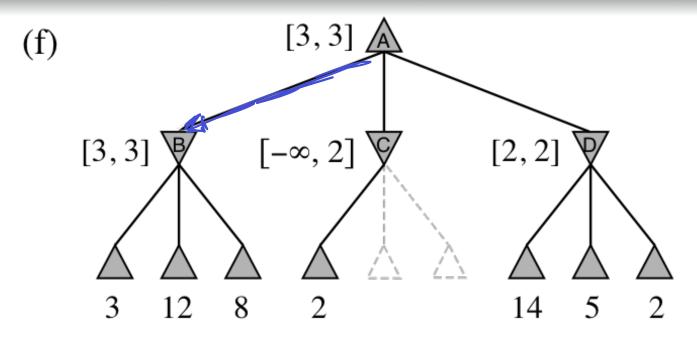




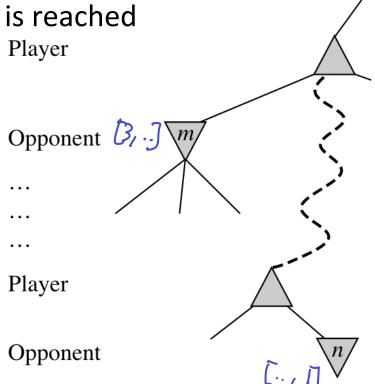








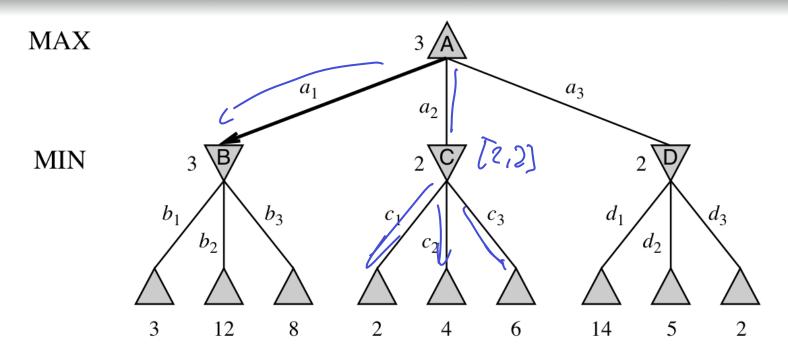
- Key idea: For each node n in minimax tree keep track of
 - α: Best value for MAX player if n is reached
 - β: Best value for MIN player if n is reached
- Never need to explore consequences of actions for which β<α
- Avoid exploring "provably suboptimal" parts
 of minimax tree



α - β -pruning algorithm

```
function ALPHA-BETA-DECISION(state) returns an action
   return the a in ACTIONS(state) maximizing MIN-VALUE(RESULT(a, state))
function Max-Value(state, \alpha, \beta) returns a utility value
   inputs: state, current state in game
             \alpha, the value of the best alternative for MAX along the path to state
             \beta, the value of the best alternative for MIN along the path to state
   if TERMINAL-Test(state) then return Utility(state)
   v \leftarrow -\infty
   for a, s in Successors(state) do
      v \leftarrow \text{Max}(v, \text{Min-Value}(s, \alpha, \beta))
      if v \geq \beta then return v
      \alpha \leftarrow \text{Max}(\alpha, v)
   return v
function MIN-Value (state, \alpha, \beta) returns a utility value
   same as MAX-VALUE but with roles of \alpha, \beta reversed
```

Does move ordering matter?



Move ordering matters a lot

- Worst case: No improvement 0(6")
- Best case (ideal ordering):

Random ordering:

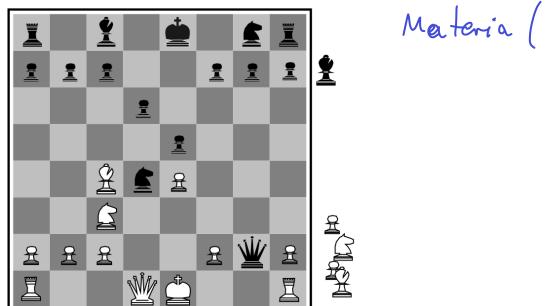
g):
$$O(6^{m/2})$$
 $Can get$ fairly close $O((b/log b)^m)$ m practice

How to find a good ordering?

Use IDS, "vanember" best moves from Challower tras

Large state spaces

- Typical branching factor in chess: 35
- Computing the complete minimax tree is intractable
- Instead: Cut off search, and replace utility(s) with eval(s)
 - eval(s) is heuristic value of state s



Materia (# black/hite
pauns, vools, ...)

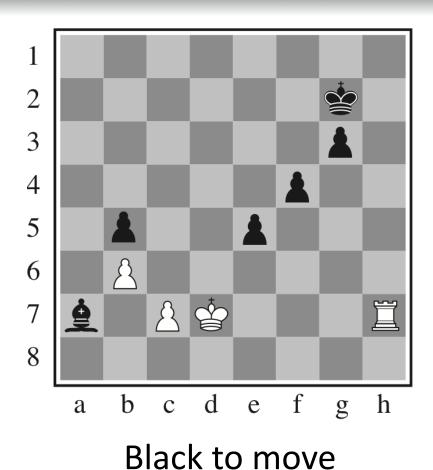
Developing evaluation functions

- This is where expert knowledge comes in
- Typical approach:
 - Select features $f_1,...,f_n$ that may be useful, e.g., value of pieces on board, positions of pieces, ...
 - Learn weights from examples

$$eval(s) = \sum_{i=1}^{n} w_i f_i(s)$$

- Deep Blue used ~6,000 different features!
- Often, reinforcement learning is very useful here (e.g., TD-gammon beats world champion in backgammon)

Problems with cutoff search



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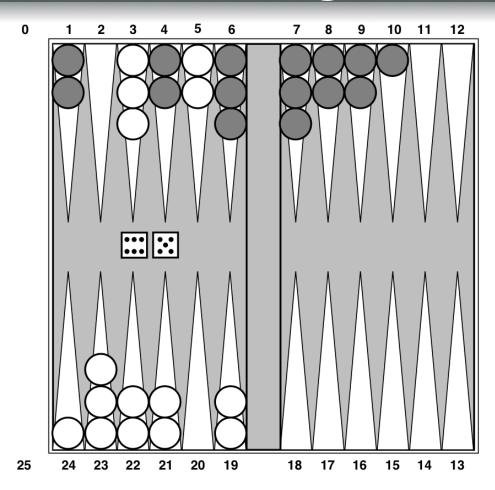
Taming the horizon effect

- Quiescence search
 - Evaluation function also evaluates "stability" (e.g., strong captures, etc.)
 - Cutoff postponed if position is unstable
 - Search time no longer constant
- Singular extension
 - Search deeper if a node's value is much better than its siblings'
 - Reduces effective branching factor
 - Can search much longer sequences (even 30-40ply)

Playing world class chess

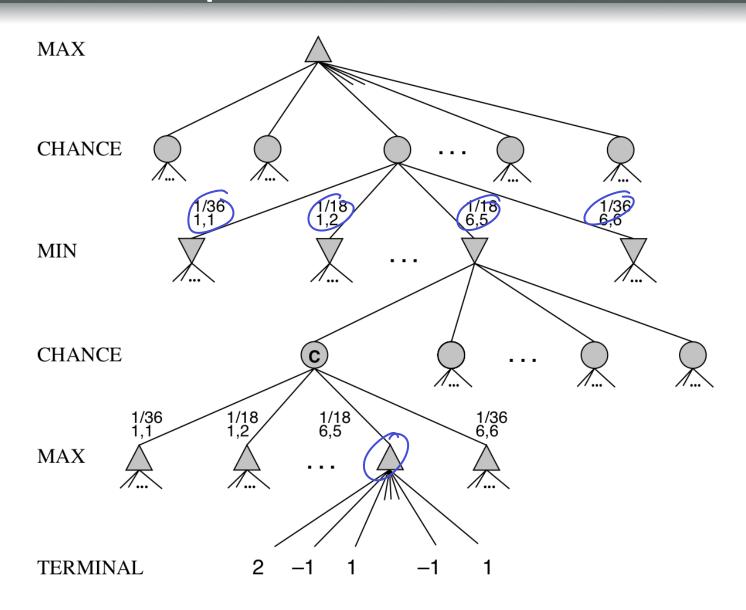
- Current PCs can evaluate ~200 million nodes / 3 min
- Minimax search: ~5 ply lookahead
- With α - β pruning: ~10 ply
- Further improvements:
 - Quiescence search: Only evaluate "stable" positions
 - Transposition tables: Remember states evaluated before
 - Singular extensions: Expand tree if there is singular best move
 - Null move heuristic: Get lower bound by letting opp. move 2x
 - Precompute endgames (all 5, some 6 piece positions)
 - Opening library (up to ~30ply in first couple moves)
- Hydra: 18 ply lookahead (on 64 processor cluster)

Stochastic games



- Two "types" of uncertainty
- Adversarial and stochastic

ExpectiMiniMax tree



Solving stochastic games

- ExpectiMiniMax used to calculate optimal move
- Defined inductively:

If *n* is terminal node (or cutoff):

• Value is utility (n.state) (or eval(n.state))

If n is MAX node:

Value is highest value of all successor node values

If *n* is MIN node

Value is lowest value of all successor node values

If *n* is CHANCE node

 Value is (weighted) average of all successor node values

Dealing with large state spaces

- Backgammon:
 - 21 possible roles with 2 die; ~20 legal moves
 - #nodes for depth 4 tree:

- As depth increases, reaching any particular node becomes exponentially unlikely
 - Lookahead becomes less valuable
 - α - β -pruning much less useful: world just won't play along!
- TD-gammon competitive with best human players:
 - Uses only 2 ply lookahead!
 - But very carefully trained evaluation function